

## Solutions to Exercises

### Chapter 2

#### 2.1 Two-oracle variant of the PAC model

- Assume that  $\mathcal{C}$  is efficiently PAC-learnable using  $\mathcal{H}$  in the standard PAC model using algorithm  $\mathcal{A}$ . Consider the distribution  $\mathcal{D} = \frac{1}{2}(\mathcal{D}_- + \mathcal{D}_+)$ . Let  $h \in \mathcal{H}$  be the hypothesis output by  $\mathcal{A}$ . Choose  $\delta$  such that:

$$\mathbb{P}[R_{\mathcal{D}}(h) \leq \epsilon/2] \geq 1 - \delta.$$

From

$$\begin{aligned} R_{\mathcal{D}}(h) &= \mathbb{P}_{x \sim \mathcal{D}}[h(x) \neq c(x)] \\ &= \frac{1}{2} \left( \mathbb{P}_{x \sim \mathcal{D}_-}[h(x) \neq c(x)] + \mathbb{P}_{x \sim \mathcal{D}_+}[h(x) \neq c(x)] \right) \\ &= \frac{1}{2}(R_{\mathcal{D}_-}(h) + R_{\mathcal{D}_+}(h)), \end{aligned}$$

it follows that:

$$\mathbb{P}[R_{\mathcal{D}_-}(h) \leq \epsilon] \geq 1 - \delta \quad \text{and} \quad \mathbb{P}[R_{\mathcal{D}_+}(h) \leq \epsilon] \geq 1 - \delta.$$

This implies two-oracle PAC-learning with the same computational complexity.

- Assume now that  $\mathcal{C}$  is efficiently PAC-learnable in the two-oracle PAC model. Thus, there exists a learning algorithm  $\mathcal{A}$  such that for  $c \in \mathcal{C}$ ,  $\epsilon > 0$ , and  $\delta > 0$ , there exist  $m_-$  and  $m_+$  polynomial in  $1/\epsilon$ ,  $1/\delta$ , and  $\text{size}(c)$ , such that if we draw  $m_-$  negative examples or more and  $m_+$  positive examples or more, with confidence  $1 - \delta$ , the hypothesis  $h$  output by  $\mathcal{A}$  verifies:

$$\mathbb{P}[R_{\mathcal{D}_-}(h) \leq \epsilon] \geq 1 - \delta \quad \text{and} \quad \mathbb{P}[R_{\mathcal{D}_+}(h) \leq \epsilon] \geq 1 - \delta.$$

Now, let  $\mathcal{D}$  be a probability distribution over negative and positive examples. If we could draw  $m$  examples according to  $\mathcal{D}$  such that  $m \geq \max\{m_-, m_+\}$ ,  $m$  polynomial in  $1/\epsilon$ ,  $1/\delta$ , and  $\text{size}(c)$ , then two-oracle PAC-learning would imply standard PAC-learning:

$$\begin{aligned} \mathbb{P}[R_{\mathcal{D}}(h)] &\leq \mathbb{P}[R_{\mathcal{D}}(h)|c(x) = 0] \mathbb{P}[c(x) = 0] + \mathbb{P}[R_{\mathcal{D}}(h)|c(x) = 1] \mathbb{P}[c(x) = 1] \\ &\leq \epsilon(\mathbb{P}[c(x) = 0] + \mathbb{P}[c(x) = 1]) = \epsilon. \end{aligned}$$

If  $\mathcal{D}$  is not too biased, that is, if the probability of drawing a positive example, or that of drawing a negative example is more than  $\epsilon$ , it is not hard to show, using Chernoff bounds or just Chebyshev's inequality, that drawing a polynomial number of examples in  $1/\epsilon$  and  $1/\delta$  suffices to guarantee that  $m \geq \max\{m_-, m_+\}$  with high confidence.

Otherwise,  $\mathcal{D}$  is biased toward negative (or positive examples), in which case returning  $h = h_0$  (respectively  $h = h_1$ ) guarantees that  $\mathbb{P}[R_{\mathcal{D}}(h)] \leq \epsilon$ .

To show the claim about the not-too-biased case, let  $S_m$  denote the number of positive examples obtained when drawing  $m$  examples when the probability of a positive example is  $\epsilon$ . By Chernoff bounds,

$$\mathbb{P}[S_m \leq (1 - \alpha)m\epsilon] \leq e^{-m\epsilon\alpha^2/2}.$$

We want to ensure that at least  $m_+$  examples are found. With  $\alpha = \frac{1}{2}$  and  $m = \frac{2m_+}{\epsilon}$ ,

$$\mathbb{P}[S_m > m_+] \leq e^{-m_+/4}.$$

Setting the bound to be less than or equal to  $\delta/2$ , leads to the following condition on  $m$ :

$$m \geq \min\left\{\frac{2m_+}{\epsilon}, \frac{8}{\epsilon} \log \frac{2}{\delta}\right\}$$

A similar analysis can be done in the case of negative examples. Thus, when  $\mathcal{D}$  is not too biased, with confidence  $1 - \delta$ , we will find at least  $m_-$  negative and  $m_+$  positive examples if we draw  $m$  examples, with

$$m \geq \min\left\{\frac{2m_+}{\epsilon}, \frac{2m_-}{\epsilon}, \frac{8}{\epsilon} \log \frac{2}{\delta}\right\}.$$

In both solutions, our training data is the set  $T$  and our learned concept  $L(T)$  is the tightest circle (with minimal radius) which is consistent with the data.

## 2.5 Triangles

As in the case of axis-aligned rectangles, consider three regions  $r_1, r_2, r_3$ , along the sides of the target concept as indicated in figure E.6. Note that the triangle formed by the points  $A'', B'', C''$  is similar to  $ABC$  (same angles) since  $A''B''$  must be parallel to  $AB$ , and similarly for the other sides.

Assume that  $\mathbb{P}[ABC] > \epsilon$ , otherwise the statement would be trivial. Consider a triangle  $A'B'C'$  similar to  $ABC$  and consistent with the training sample and such that it meets all three regions  $r_1, r_2, r_3$ .

Since it meets  $r_1$ , the line  $A'B'$  must be below  $A''B''$ . Since it meets  $r_2$  and  $r_3$ ,  $A'$  must be in  $r_2$  and  $B'$  in  $r_3$  (see figure E.6). Now, since the angle  $A'B'C'$  is equal to  $A''B''C''$ ,  $C'$  must be necessarily above  $C''$ . This implies that triangle  $A'B'C'$  contains  $A''B''C''$ , and thus  $\text{error}(A'B'C') \leq \epsilon$ .

$$\text{error}(A'B'C') > \epsilon \implies \exists i \in \{1, 2, 3\} : A'B'C' \cap r_i = \emptyset.$$

Thus, by the union bound,

$$\mathbb{P}[\text{error}(A'B'C') > \epsilon] \leq \sum_{i=1}^3 \mathbb{P}[A'B'C' \cap r_i = \emptyset] \leq 3(1 - \epsilon/3)^m \leq 3e^{-3m\epsilon}.$$

Setting  $\delta$  to match the right-hand side gives the sample complexity  $m \geq \frac{3}{\epsilon} \log \frac{3}{\delta}$ .

## Chapter 3

### 3.3 Growth function of linear combinations

- (a)  $\{\mathcal{X}^+ \cup \{\mathbf{x}_{m+1}\}, \mathcal{X}^-\}$  and  $\{\mathcal{X}^+, \mathcal{X}^- \cup \{\mathbf{x}_{m+1}\}\}$  are linearly separable by a hyperplane going through the origin if and only if there exists  $\mathbf{w}_1 \in \mathbb{R}^d$  such that

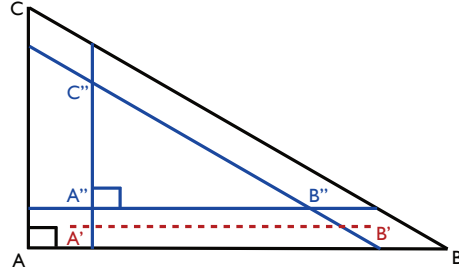
$$\forall \mathbf{x} \in \mathcal{X}^+, \mathbf{w}_1 \cdot \mathbf{x} > 0 \quad \forall \mathbf{x} \in \mathcal{X}^-, \mathbf{w}_1 \cdot \mathbf{x} < 0, \text{ and } \mathbf{w}_1 \cdot \mathbf{x}_{m+1} > 0 \quad (\text{E.35})$$

and there exists  $\mathbf{w}_2 \in \mathbb{R}^d$  such that

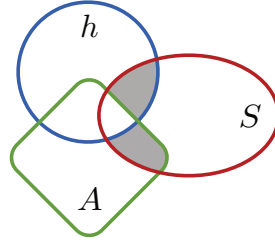
$$\forall \mathbf{x} \in \mathcal{X}^+, \mathbf{w}_2 \cdot \mathbf{x} > 0 \quad \forall \mathbf{x} \in \mathcal{X}^-, \mathbf{w}_2 \cdot \mathbf{x} < 0, \text{ and } \mathbf{w}_2 \cdot \mathbf{x}_{m+1} < 0. \quad (\text{E.36})$$

For any  $\mathbf{w}_1, \mathbf{w}_2$ , the function  $f: (t \mapsto t\mathbf{w}_1 + (1-t)\mathbf{w}_2) \cdot \mathbf{x}_{m+1}$  is continuous over  $[0, 1]$ . (E.44) and (E.45) hold iff  $f(0) < 0$  and  $f(1) > 0$ , that is iff there exists  $\mathbf{w} = t_0\mathbf{w}_1 + (1-t_0)\mathbf{w}_2$  linearly separating  $\{\mathcal{X}^+, \mathcal{X}^-\}$  and such at  $\mathbf{w} \cdot \mathbf{x}_{m+1} = 0$ .

- (b) Repeating the formula, we obtain  $C(m, d) = \sum_{k=0}^{m-1} \binom{m-1}{k} C(1, d-k)$ . Since,  $C(1, n) = 2$  if  $n \geq 1$  and  $C(1, n) = 0$  otherwise, the result follows.



**Figure E.5**  
Rectangle triangles.



**Figure E.6**  
Illustration of  $(h\Delta A) \cap S = (h \cap S)\Delta(A \cap S)$  shown in gray.

(c) This is a direct application of the result of the previous question.

### 3.25 VC-dimension of symmetric difference of concepts

Fix a set  $S$ . We can show that the number of classifications of  $S$  using  $\mathcal{H}$  is the same as when using  $\mathcal{H}\Delta A$ . The set of classifications obtained using  $\mathcal{H}$  can be identified with  $\{S \cap h : h \in \mathcal{H}\}$  and the set of classifications using  $\mathcal{H}\Delta A$  can be identified with  $\{S \cap (h\Delta A) : h \in \mathcal{H}\}$ . Observe that for any  $h \in \mathcal{H}$ ,

$$S \cap (h\Delta A) = (S \cap h)\Delta(S \cap A). \quad (\text{E.37})$$

Figure E.7 helps illustrate this equality in a special case. Now, in view of this inequality, if  $S \cap (h\Delta A) = S \cap (h'\Delta A)$  for  $h, h' \in \mathcal{H}$ , then

$$(S \cap h)\Delta \mathcal{B} = (S \cap h')\Delta \mathcal{B}, \quad (\text{E.38})$$

with  $\mathcal{B} = S \cap A$ . Since two sets that have the same symmetric differences with respect to a set  $\mathcal{B}$  must be equal, this implies

$$S \cap h = S \cap h'. \quad (\text{E.39})$$

This shows that  $\phi$  defined by

$$\begin{aligned} \phi: S \cap \mathcal{H} &\rightarrow S \cap (\mathcal{H}\Delta A) \\ S \cap h &\mapsto S \cap (h\Delta A) \end{aligned}$$

is a bijection, and thus that the sets  $S \cap \mathcal{H}$  and  $S \cap (\mathcal{H}\Delta A)$  have the same cardinality.

## Chapter 5

### 5.3 Importance weighted SVM

The modified primal optimization problem can be written as

$$\begin{aligned} & \text{minimize} && \frac{1}{2} \|w\|^2 + C \sum_{i=1}^m \xi_i p_i \\ & \text{subject to} && y_i [w \cdot x_i + b] \geq 1 - \xi_i. \end{aligned}$$

The Lagrangian holding for all  $w, b, \alpha_i \geq 0, \beta_i \geq 0$  is then

$$\begin{aligned} L(w, b, \alpha) = & \frac{1}{2} \|w\|^2 + C \sum_{i=1}^m \xi_i p_i \\ & - \sum_{i=1}^m \alpha_i [y_i (w \cdot x_i + b) - 1 + \xi_i] - \sum_{i=1}^m \beta_i \xi_i. \end{aligned} \quad (\text{E.40})$$

Then  $\frac{\partial L}{\partial w}$  and  $\frac{\partial L}{\partial b}$  are the same as for the regular non-separable SVM optimization problem. We also have  $\frac{\partial L}{\partial \xi_i} = Cp_i - \alpha_i - \beta_i$ . Thus, to satisfy the KKT conditions we have for all  $i \in [m]$ ,

$$w = \sum_{i=1}^m \alpha_i y_i x_i \quad (\text{E.41})$$

$$\sum_{i=1}^m \alpha_i y_i = 0 \quad (\text{E.42})$$

$$\alpha_i + \beta_i = Cp_i \quad (\text{E.43})$$

$$\alpha_i [y_i (w \cdot x_i + b) - 1 + \xi_i] = 0 \quad (\text{E.44})$$

$$\beta_i \xi_i = 0. \quad (\text{E.45})$$

Plugging equation E.79 into equation E.78, we get

$$\begin{aligned} L = & \frac{1}{2} \left\| \sum_{i=1}^m \alpha_i y_i x_i \right\|^2 + C \sum_{i=1}^m \xi_i p_i - \sum_{i,j=1}^m \alpha_i \alpha_j y_i y_j (x_i \cdot x_j) \\ & - \sum_{i=1}^m \alpha_i y_i b + \sum_{i=1}^m \alpha_i - \sum_{i=1}^m \alpha_i \xi_i - \sum_{i=1}^m \beta_i \xi_i. \end{aligned} \quad (\text{E.46})$$

Using equation E.81, we can simplify:

$$L = \sum_{i=1}^m \alpha_i - \frac{1}{2} \left\| \sum_{i=1}^m \alpha_i y_i x_i \right\|^2,$$

meaning that the objective function is the same as in the regular SVM problem. The difference is in the constraints on the optimization. Recall that our dual form holds for  $\beta_i \geq 0$ . Using again equation E.81, our optimization problem is to maximize  $L$  subject to the constraints:

$$\forall i \in [m], 0 \leq \alpha_i \leq Cp_i \wedge \sum_{i=1}^m \alpha_i y_i = 0.$$

### 5.6 Sparse SVM

(a) Let

$$\mathbf{x}'_i = (y_1 \mathbf{x}_i \cdot \mathbf{x}_1, \dots, y_m \mathbf{x}_i \cdot \mathbf{x}_m).$$

Then the optimization problem becomes

$$\begin{aligned} \min_{\boldsymbol{\alpha}, b, \xi} \quad & \frac{1}{2} \|\boldsymbol{\alpha}\|^2 + C \sum_{i=1}^m \xi_i \\ \text{subject to} \quad & y_i (\boldsymbol{\alpha} \cdot \mathbf{x}'_i + b) \geq 1 - \xi_i \\ & \xi_i, \alpha_i \geq 0, i \in [m], \end{aligned}$$

which is the standard formulation of the primal SVM optimization problem on samples  $\mathbf{x}'_i$ , modulo the non-negativity constraints on  $\alpha_i$ .

(b) The Lagrangian of (1) for all  $\alpha_i \geq 0, \xi_i \geq 0, b, \alpha'_i \geq 0, \beta_i \geq 0, \gamma_i \geq 0, i \in [m]$  is

$$L = \frac{1}{2} \|\boldsymbol{\alpha}\|^2 + C \sum_{i=1}^m \xi_i - \sum_{i=1}^m \alpha'_i (y_i (\boldsymbol{\alpha} \cdot \mathbf{x}'_i + b) - 1 + \xi_i) - \sum_{i=1}^m \beta_i \xi_i - \sum_{i=1}^m \gamma_i \alpha_i,$$

and the KKT conditions are

$$\nabla_{\boldsymbol{\alpha}} L = 0 \quad \Leftrightarrow \quad \boldsymbol{\alpha} = \sum_{i=1}^m \alpha'_i y_i \mathbf{x}'_i + \boldsymbol{\gamma}$$

$$\nabla_b L = 0 \quad \Leftrightarrow \quad \sum_{i=1}^m \alpha'_i y_i = 0$$

$$\nabla_{\xi_i} L = 0 \quad \Leftrightarrow \quad \alpha'_i + \beta_i = C$$

and

$$\begin{aligned} \alpha'_i (y_i (\boldsymbol{\alpha} \cdot \mathbf{x}'_i + b) - 1 + \xi_i) &= 0 \\ \beta_i \xi_i &= 0 \\ \gamma_i \alpha_i &= 0. \end{aligned}$$

Using the KKT conditions on  $L$  we get

$$\begin{aligned} L &= \frac{1}{2} \left( \sum_{i=1}^m \alpha'_i y_i \mathbf{x}'_i + \boldsymbol{\gamma} \right) \cdot \left( \sum_{j=1}^m \alpha'_j y_j \mathbf{x}'_j + \boldsymbol{\gamma} \right) + C \sum_{i=1}^m \xi_i \\ &\quad - \sum_{i=1}^m \alpha'_i \left( y_i \left( \left( \sum_{j=1}^m \alpha'_j y_j \mathbf{x}'_j + \boldsymbol{\gamma} \right) \cdot \mathbf{x}'_i + b \right) - 1 + \xi_i \right) \\ &\quad - \sum_{i=1}^m \beta_i \xi_i - \sum_{i=1}^m \gamma_i \alpha_i \\ &= -\frac{1}{2} \sum_{i=1}^m \alpha'_i y_i \mathbf{x}'_i \cdot \left( \sum_{j=1}^m \alpha'_j y_j \mathbf{x}'_j + \boldsymbol{\gamma} \right) + \frac{1}{2} \boldsymbol{\gamma} \cdot \boldsymbol{\gamma} \\ &\quad + \sum_{i=1}^m C \xi_i - \alpha'_i (y_i b - 1 + \xi_i) \\ &= \sum_{i=1}^m \alpha'_i - \frac{1}{2} \sum_{i,j=1}^m \alpha'_i \alpha'_j y_i y_j \mathbf{x}'_i{}^\top (\mathbf{x}'_j + \boldsymbol{\gamma}) \\ &\quad + \sum_{i=1}^m (C - \alpha'_i) \xi_i - \sum_{i=1}^m \alpha'_i y_i b. \end{aligned}$$

Thus the dual optimization problem is

$$\begin{aligned} \max_{\alpha', \gamma} \quad & \sum_{i=1}^m \alpha'_i - \frac{1}{2} \sum_{i,j=1}^m \alpha'_i \alpha'_j y_i y_j \mathbf{x}'_i \cdot (\mathbf{x}'_j + \gamma) \\ \text{subject to} \quad & \sum_{i=1}^m \alpha'_i y_i = 0 \\ & 0 \leq \alpha'_i \leq C, \gamma_i \geq 0, i \in [m]. \end{aligned}$$

## Chapter 6

### 6.18 Metrics and kernels

- (a) If  $K$  is an NDS kernel, then by theorem 6.16 the kernel  $K'$  defined for any  $x_0 \in \mathcal{X}$  by:

$$K'(x, x') = \frac{1}{2} [K(x, x_0) + K(x', x_0) - K(x, x')]$$

is a PDS kernel ( $K(x_0, x_0) = 0$ ). Let  $\mathbb{H}$  be the reproducing Hilbert space associated to  $K'$ . There exists a mapping  $\Phi(x)$  from  $\mathcal{X}$  to  $\mathbb{H}$  such that  $\forall x, x' \in \mathcal{X}, K'(x, x') = \Phi(x) \cdot \Phi(x')$ . Then,

$$\begin{aligned} \|\Phi(x) - \Phi(x')\|^2 &= K'(x, x) + K'(x', x') - 2K'(x, x') \\ &= \frac{1}{2} [2K(x, x_0) - K(x, x)] + \\ &\quad \frac{1}{2} [2K(x', x_0) - K(x', x')] - \\ &\quad [K(x, x_0) + K(x', x_0) - K(x, x')] \\ &= K(x, x') \end{aligned}$$

It is then straightforward to show that  $\sqrt{K}$  is a metric.

- (b) Suppose that  $K(x, x') = \exp(-|x - x'|^p)$ ,  $x, x' \in \mathbb{R}$ , is positive definite for  $p > 2$ . Then, for any  $t > 0$ ,  $\{x_1, \dots, x_n\} \subseteq X$ ,  $\{c_1, \dots, c_n\} \subseteq \mathbb{R}$ ,

$$\sum_{i,j=1}^n c_i c_j \exp(-t|x_j - x_k|^p) = \sum_{i,j=1}^n c_i c_j \exp(-|t^{1/p} x_j - t^{1/p} x_k|^p) \geq 0$$

Thus, by theorem 6.17,  $K'(x, x') = |x - x'|^p$  is an NDS kernel. But,  $\sqrt{K'}$  is not a metric for  $p > 2$  since it does not verify the triangle inequality (take  $x = 1$ ,  $x' = 2$ ,  $x'' = 3$ ), which contradicts part (a).

- (c) If  $a < 0$  or  $b < 0$ ,  $a\|x\|^2 + b < 0$  for some non-null vectors  $x$ . For such values,  $K(x, x) = \tanh(a\|x\|^2 + b) < 0$ . The kernel is thus not PDS and the SVM training may not converge to an optimal value. The equivalent neural network may also converge to a local minimum.

### 6.19 Sequence kernels

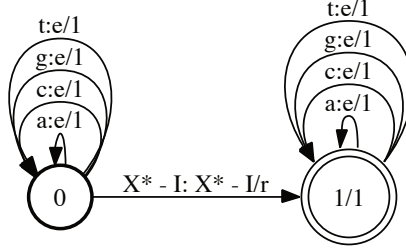
- (a)  $X^* - I$  is a regular language and can be represented by a finite automaton.  $K$  can thus be defined by

$$\forall x, y \in X^*, \quad K(x, y) = [[T \circ T^{-1}]](x, y), \quad (\text{E.47})$$

where  $T$  is the weighted transducer shown in figure E.14. Thus,  $K$  is a rational kernel and in view of the theorem 6.21, it is positive definite symmetric.

- (b) Let  $M_{X^* - I}$  be the minimal automaton representing  $X^* - I$ . The transducer  $T$  of figure E.14 can be constructed using  $M_{X^* - I}$ . Then,  $|T| = |M_{X^* - I}| + 8$ . Using composition of weighted transducers, the running time complexity of the computation of the algorithm is:

$$O(|x||y||T \circ T^{-1}|) = O(|x||y||T|^2) = O(|x||y||M_{X^* - I}|^2). \quad (\text{E.48})$$

**Figure E.7**

Weighted transducer  $T$ .  $e$  represents the empty string, and  $r = \rho$ .  $X^* - I$  stands for a finite automaton accepting  $X^* - I$ .

- (c) The set of strings  $Y$  over the alphabet  $X$  of length less than  $n$  form a regular language since they can be described by:

$$Y = \bigcup_{i=0}^{n-1} X^i. \quad (\text{E.49})$$

Thus,  $Y_1 = Y \cap (X^* - I)$  and  $Y_2 = (X^* - I) - Y_1$  are also regular languages. It suffices to replace in the transducer  $T$  of figure E.14 the transition labeled with  $X^* - I : X^* - I/\rho$  with two transitions:

- $Y_1 : Y_1/\rho_1$ , and
- $Y_2 : Y_2/\rho_2$ ,

with the same origin and destination states and with  $Y_1$  and  $Y_2$  denoting finite automata representing them. The kernel is thus still rational and PDS since it is of the form  $T' \circ T'^{-1}$ .

## Chapter 7

### 7.8 Simplified AdaBoost

- (a) As in the standard case, we can show that

$$\hat{R}(h) \leq \prod_{t=1}^T Z_t, \quad (\text{E.50})$$

and that

$$Z_t = (1 - \epsilon_t)e^{-\alpha} + \epsilon_t e^{\alpha}. \quad (\text{E.51})$$

By definition of  $\gamma$  and the fact that  $e^{\alpha} - e^{-\alpha} > 0$  for all  $\alpha > 0$ ,

$$Z_t = \epsilon_t(e^{\alpha} - e^{-\alpha}) + e^{-\alpha} \quad (\text{E.52})$$

$$\leq (1 - \gamma)(e^{\alpha} - e^{-\alpha}) + e^{-\alpha} \quad (\text{E.53})$$

$$= \left(\frac{1}{2} - \gamma\right)e^{\alpha} + \left(\frac{1}{2} + \gamma\right)e^{-\alpha} = u(\alpha). \quad (\text{E.54})$$

$u(\alpha)$  is minimized for

$$\left(\frac{1}{2} - \gamma\right)e^{\alpha} = \left(\frac{1}{2} + \gamma\right)e^{-\alpha}, \quad (\text{E.55})$$

that is, for

$$\alpha = \frac{1}{2} \log \frac{\frac{1}{2} + \gamma}{\frac{1}{2} - \gamma}. \quad (\text{E.56})$$

Tighter bounds on the product of the  $Z_t$ s can lead to better values for  $\alpha$ .

- (b) As in the standard case, at round  $t$ , the probability mass assigned to correctly classified points is  $p_+ = (1 - \epsilon_t)e^{-\alpha}$  and the probability mass assigned to the misclassified points is  $p_- = \epsilon_t e^\alpha$ . Thus,

$$\frac{p_-}{p_+} = \frac{\epsilon_t}{1 - \epsilon_t} \frac{\frac{1}{2} + \gamma}{\frac{1}{2} - \gamma} \leq \frac{\frac{1}{2} - \gamma}{\frac{1}{2} + \gamma} \frac{\frac{1}{2} + \gamma}{\frac{1}{2} - \gamma} = 1. \quad (\text{E.57})$$

This contrasts with AdaBoost's property.

- (c)

$$Z_t \leq \left(\frac{1}{2} - \gamma\right)e^\alpha + \left(\frac{1}{2} + \gamma\right)e^{-\alpha} \quad (\text{E.58})$$

$$= \left(\frac{1}{2} - \gamma\right)\sqrt{\frac{\frac{1}{2} + \gamma}{\frac{1}{2} - \gamma}} + \left(\frac{1}{2} + \gamma\right)\sqrt{\frac{\frac{1}{2} - \gamma}{\frac{1}{2} + \gamma}} \quad (\text{E.59})$$

$$= 2\sqrt{\left(\frac{1}{2} + \gamma\right)\left(\frac{1}{2} - \gamma\right)}. \quad (\text{E.60})$$

Thus, the empirical error can be bounded as follows:

$$\widehat{R}_S(h) \leq \prod_{t=1}^T Z_t \quad (\text{E.61})$$

$$\leq [2\sqrt{\left(\frac{1}{2} + \gamma\right)\left(\frac{1}{2} - \gamma\right)}]^T \quad (\text{E.62})$$

$$= (1 - 4\gamma^2)^{T/2} \quad (\text{E.63})$$

$$\leq e^{-2\gamma^2 T}. \quad (\text{E.64})$$

- (d) If  $\widehat{R}_S(h) = \frac{1}{m} \sum_{i=1}^m 1_{y_i f(x_i) \leq 0} \leq \frac{1}{m}$ , then clearly  $\widehat{R}_S(h) = 0$ . Using the bound obtained in the previous question, if  $e^{-2\gamma^2 T} < \frac{1}{m}$ , the empirical error is zero. This can be rewritten as

$$T > \frac{\log m}{2\gamma^2}. \quad (\text{E.65})$$

- (e) Using the bound for the consistent case,

$$\mathbb{P}[R(h) > \epsilon] \leq 2\Pi_e(2m)2^{-\frac{m\epsilon}{2}} \leq 2\left(\frac{2em}{d}\right)^d 2^{-\frac{m\epsilon}{2}}. \quad (\text{E.66})$$

Setting the right-hand side to  $\delta$ , with probability at least  $1 - \delta$ , the following bound holds for that consistent hypothesis:

$$\text{error}_{\mathcal{D}}(\mathcal{H}) \leq \frac{2}{m} \left( d \log_2 \frac{2em}{d} + \log_2 \frac{2}{\delta} \right), \quad (\text{E.67})$$

with  $d = 2(s+1)T \log_2(eT)$  and  $T = \left\lfloor \frac{\log m}{2\gamma^2} \right\rfloor + 1$ .

The bound is vacuous for  $\gamma(m) = O(\sqrt{\frac{\log m}{m}})$ . This could suggest overfitting.

## 7.11 HingeBoost

- (a) Since the hinge loss is convex, its composition with affine function of  $\alpha$  is also convex and  $F$  is convex as sum of convex functions.

For the existence of one-sided directional derivatives, one can use the fact that any convex function has one-sided directional derivatives or alternatively, that our specific function is the sum of piecewise affine functions, which are also known to have one-sided directional derivatives (think of one-dimensional hinge loss).



- (b) Distinguishing different cases depending on the value of  $y_i f_{t-1}(x_i) = 1$ , it is straightforward to derive the following expressions for all  $j \in [N]$ :

$$F'_+(\alpha_{t-1}, e_j) = \sum_{i=1}^m -y_i h_j(x_i) [1_{y_i f_{t-1}(x_i) < 1} + 1_{(y_i h_j(x_i) < 0) \wedge (y_i f_{t-1}(x_i) = 1)}]$$

$$F'_-(\alpha_{t-1}, e_j) = \sum_{i=1}^m -y_i h_j(x_i) [1_{y_i f_{t-1}(x_i) < 1} + 1_{(y_i h_j(x_i) > 0) \wedge (y_i f_{t-1}(x_i) = 1)}].$$

The key here is that when  $y_i f_{t-1}(x_i) \neq 1$ , each term in the sum will be either 0 or the affine function independent of  $y_i h_j(x_i)$ . On the other hand, when  $y_i f_{t-1}(x_i) = 1$ , the sign of  $y_i h_j(x_i)$  determines whether the finite differences will extend into the 0 portion of the affine portion of the term.

---

- (c)

```

HINGEBOOST( $S = ((x_1, y_1), \dots, (x_m, y_m))$ )
1   $f \leftarrow 0$ 
2  for  $j \leftarrow 1$  to  $N$  do
3       $r \leftarrow \sum_{i=1}^m -y_i h_j(x_i) [1_{y_i f(x_i) < 1} + 1_{(y_i h_j(x_i) < 0) \wedge (y_i f(x_i) = 1)}]$ 
4       $l \leftarrow \sum_{i=1}^m -y_i h_j(x_i) [1_{y_i f(x_i) < 1} + 1_{(y_i h_j(x_i) > 0) \wedge (y_i f(x_i) = 1)}]$ 
5      if  $(l \leq 0) \wedge (r \geq 0)$  then
6           $d[j] \leftarrow 0$ 
7      elseif  $(l \leq r)$  then
8           $d[j] \leftarrow r$ 
9      else  $d[j] \leftarrow l$ 
10 for  $t \leftarrow 1$  to  $T$  do
11      $k \leftarrow \operatorname{argmin}_{j \in [N]} |d[j]|$ 
12      $\eta \leftarrow \operatorname{argmin}_{\eta \geq 0} G(f + \eta h_k) \triangleright \text{line search}$ 
13      $f \leftarrow f + \eta h_k$ 
14 return  $f$ 

```

---

## Chapter 8

### 8.2 Generalized mistake bound

The bound is unaffected, as shown by the following, using the same definitions and steps as in this chapter:

$$\begin{aligned}
M\rho &\leq \frac{\mathbf{v} \cdot \sum_{t \in I} y_t \mathbf{x}_t}{\|\mathbf{v}\|} \\
&= \frac{\mathbf{v} \cdot \sum_{t \in I} (\mathbf{w}_{t+1} - \mathbf{w}_t)/\eta}{\|\mathbf{v}\|} \quad (\text{definition of updates}) \\
&= \frac{\mathbf{v} \cdot \mathbf{w}_{T+1}}{\eta \|\mathbf{v}\|} \\
&\leq \|\mathbf{w}_{T+1}\|/\eta \quad (\text{Cauchy-Schwarz ineq.}) \\
&= \|\mathbf{w}_{t_m} + \eta y_{t_m} \mathbf{x}_{t_m}\|/\eta \quad (t_m \text{ largest } t \text{ in } I) \\
&= \left[ \|\mathbf{w}_{t_m}\|^2 + \eta^2 \|\mathbf{x}_{t_m}\|^2 + \underbrace{2\eta y_{t_m} \mathbf{w}_{t_m} \cdot \mathbf{x}_{t_m}}_{\leq 0} \right]^{1/2}/\eta \\
&\leq \left[ \|\mathbf{w}_{t_m}\|^2 + \eta^2 R^2 \right]^{1/2}/\eta \\
&\leq \left[ M\eta^2 R^2 \right]^{1/2}/\eta = \sqrt{M}R. \quad (\text{applying the same to previous } ts \text{ in } I).
\end{aligned}$$

#### 8.10 On-line to batch — non-convex loss

(a) We use the following series of inequalities:

$$\begin{aligned}
&\min_{i \in [T]} (R(h_i) + 2c_\delta(T - i + 1)) \\
&\leq \frac{1}{T} \sum_{i=1}^T (R(h_i) + 2c_\delta(T - i + 1)) \\
&= \frac{1}{T} \sum_{i=1}^T R(h_{i-1}) + \frac{2}{T} \sum_{i=0}^{T-1} \sqrt{\frac{1}{2(T-i)} \log \frac{T(T+1)}{\delta}} \\
&< \frac{1}{T} \sum_{i=1}^T R(h_{i-1}) + \frac{2}{T} \sum_{i=0}^{T-1} \sqrt{\frac{1}{2(T-i)} \log \left( \frac{(T+1)}{\delta} \right)^2} \\
&= \frac{1}{T} \sum_{i=1}^T R(h_{i-1}) + \frac{2}{T} \sum_{i=0}^{T-1} \sqrt{\frac{1}{(T-i)} \log \frac{(T+1)}{\delta}} \\
&\leq \frac{1}{T} \sum_{i=1}^T R(h_{i-1}) + 4\sqrt{\frac{1}{T} \log \frac{T+1}{\delta}}.
\end{aligned}$$

The first inequality follows, since the minimum is always less than or equal to the average and the final inequality follows from  $\sum_{i=0}^{T-1} \sqrt{1/(T-i)} = \sum_{i=1}^T \sqrt{1/i} \leq 2\sqrt{T}$ .

(b) Coupling the inequality of part (a) with the high probability statement of lemma 8.14 to bound  $\frac{1}{T} \sum_{i=1}^T R(h_i)$  shows the desired bound.

(c) The square-root terms in part (b) can be bounded further by  $6\sqrt{\frac{1}{T} \log \frac{2(T+1)}{\delta}}$ .

Now, note that for two events  $A$  and  $B$  that each occur with probability at least  $1 - \delta$ ,

$$\begin{aligned}
\mathbb{P}[\neg A \cup \neg B] &\leq \mathbb{P}[\neg A] + \mathbb{P}[\neg B] \leq 2\delta \\
&\iff \mathbb{P}[A \wedge B] \geq 1 - 2\delta.
\end{aligned}$$

Thus, the probability that both bounds in (b) and (c) hold simultaneously is at least  $1 - 2\delta$ ; substituting  $\delta$  with  $\delta/2$  everywhere completes the bound.

## Chapter 9

9.5 Decision trees. A binary decision tree with  $n$  nodes has exactly  $n+1$  leaves. Each node can be labeled with an integer from  $\{1, \dots, N\}$  indicating which dimension is queried to make a binary split and each leaf can be labeled with  $\pm 1$  to indicate the classification made at that leaf. Fix an ordering of the nodes and leaves and consider all possible labelings of this sequence. There can be no more than  $(N+2)^{2n+1}$  distinct binary trees and, thus, the VC-dimension of this finite set of hypotheses can be no larger than  $(2n+1)\log(N+2) = O(n \log N)$ .

## Chapter 11

### 11.1 Pseudo-dimension and monotonic functions

If for some  $m > 0$ , there exists  $(t_1, \dots, t_m)$  and a set of points  $(x_1, \dots, x_m)$  that  $\mathcal{H}$  shatters, then  $\phi \circ \mathcal{H}$  can also shatter it. To see that, note that if for some  $h \in \mathcal{H}$ ,

$$h(x_i) \geq t_i,$$

then by the monotonic property of  $\phi$ ,

$$\phi(h(x_i)) \geq \phi(t_i).$$

A similar argument holds for the case  $h(x_i) < t_i$ . Thus,  $\phi \circ \mathcal{H}$  can shatter the set of points  $(x_1, \dots, x_m)$  with thresholds  $(\phi(t_1), \dots, \phi(t_m))$ , and this proves that  $\text{Pdim}(\phi \circ \mathcal{H}) \geq \text{Pdim}(\mathcal{H})$ . Since  $\phi$  is strictly monotonic, it is invertible, and a similar argument with  $\phi^{-1}$  can be used to show  $\text{Pdim}(\mathcal{H}) \geq \text{Pdim}(\phi \circ \mathcal{H})$ .

### 11.8 Optimal kernel matrix

- (a) Using the closed-form solution for the inner maximization problem  $\alpha = (\mathbf{K} + \lambda \mathbf{I})^{-1} \mathbf{y}$ , simplifies the joint optimization to a simpler minimization:

$$\min_{\mathbf{K} \succeq 0} \mathbf{y}^\top (\mathbf{K} + \lambda \mathbf{I})^{-1} \mathbf{y}, \quad \text{s.t. } \|\mathbf{K}\|_2 \leq 1.$$

Note that for any invertible matrix  $\mathbf{A}$ ,  $\mathbf{y}^\top \mathbf{A}^{-1} \mathbf{y} \geq \|\mathbf{y}\|^2 \lambda_{\min}(\mathbf{A}^{-1}) = \|\mathbf{y}\|^2 \lambda_{\max}(\mathbf{A})^{-1}$ . Thus, it is easy to see that  $\min_{\mathbf{K} \succeq 0} \mathbf{y}^\top (\mathbf{K} + \lambda \mathbf{I})^{-1} \mathbf{y} \geq \frac{\|\mathbf{y}\|^2}{1+\lambda}$  since  $\|\mathbf{K}\|_2 = \lambda_{\max}(\mathbf{K}) \leq 1$ .

We now show  $\mathbf{K} = \frac{1}{\|\mathbf{y}\|^2} \mathbf{y} \mathbf{y}^\top$  achieves this lower bound. First, note that  $(\frac{1}{\|\mathbf{y}\|^2} \mathbf{y} \mathbf{y}^\top + \lambda \mathbf{I}) \mathbf{y} = (1+\lambda) \mathbf{y}$ , so  $\mathbf{y}$  is an eigenvector of the matrix with eigenvalue  $(1+\lambda)$ . Since the matrix is invertible, it can be shown that  $\mathbf{y}$  is also an eigenvector of  $(\frac{1}{\|\mathbf{y}\|^2} \mathbf{y} \mathbf{y}^\top + \lambda \mathbf{I})^{-1}$  with eigenvalue  $\frac{1}{1+\lambda}$  (for example, consider the eigen decomposition of the matrix).

- (b) The kernel matrix alone is not useful for classifying future unseen points  $x$ , which requires computing  $\sum_{i=1}^m K(x_i, x)$  and needs access to an underlying kernel *function* that is consistent with the kernel matrix. Finding such a kernel function may be difficult in general, and furthermore the choice of function may not be unique.

## Chapter 14

### 14.1 Tighter stability bounds

- (a) No, even as  $\beta \rightarrow 0$  the generalization bound of theorem 14.2 only guarantees  $R(h_S) - \hat{R}_S(h_S) \leq M \sqrt{\frac{\log \frac{1}{\delta}}{2m}} = O(1/\sqrt{m})$ .
- (b) In this case,  $M = C/\sqrt{m}$  and  $M \sqrt{\frac{\log \frac{1}{\delta}}{2m}} = O(1/m)$ ; thus, it would suffice to have  $\beta = O(1/m^{3/2})$  in order to guarantee an  $O(1/m)$  generalization bound.

## 14.2 Quadratic hinge loss stability

We first show that the loss function is  $\sigma$ -admissible. Consider three cases:

- Both  $h(x)$  and  $h'(x)$  are correct with margin greater than 1, then

$$|L(h(x), y) - L(h'(x), y)| = 0.$$

- Only one hypothesis is correct with large enough margin. Without loss of generality assume  $h'(x)$  is correct, then

$$\begin{aligned} |L(h(x), y) - L(h'(x), y)| &= (1 - h(x)y)^2 \\ &\leq ((1 - h(x)y) - (1 - h'(x)y))^2 = (h'(x) - h(x))^2 \\ &\leq 4\sqrt{M}|h(x) - h'(x)|. \end{aligned}$$

The first inequality follows from the assumption  $1 - h'(x)y \leq 0$ , and the second inequality follows from the bounded loss assumption, which implies  $\forall h \in \mathcal{H}, |h(x)| \leq \sqrt{M} + 1 \leq 2M$ .

- Finally, we consider the case where both  $h(x)$  and  $h'(x)$  incur a loss. Without loss of generality assume  $(1 - h(x)y) \geq (1 - h'(x)y)$ , then

$$\begin{aligned} |L(h(x), y) - L(h'(x), y)| &= (1 - h(x)y)^2 - (1 - h'(x)y)^2 \\ &= ((1 - h(x)y) + (1 - h'(x)y))((1 - h(x)y) - (1 - h'(x)y)) \\ &\leq |2 - y(h(x) + h'(x))||y(h(x) - h'(x))| \leq 6\sqrt{M}|h(x) - h'(x)|. \end{aligned}$$

Thus, the quadratic hinge loss is  $\sigma$ -admissible with  $\sigma = 6\sqrt{M}$ . By proposition 14.4, SVM with quadratic hinge loss is stable with  $\beta = \frac{36r^2M}{m\lambda}$ , and using theorem 14.2 gives the following bound:

$$R(h_S) \leq \widehat{R}_S(h_S) + \frac{36r^2M}{m\lambda} + \left(\frac{72r^2M}{\lambda} + M\right)\sqrt{\frac{\log \frac{1}{\delta}}{m}}.$$

## Chapter 15

## 15.2 Double centering

- Observe that  $\|\mathbf{x}_i - \mathbf{x}_j\|^2 = (\mathbf{x}_i - \mathbf{x}_j)^\top (\mathbf{x}_i - \mathbf{x}_j) = \mathbf{x}_i^\top \mathbf{x}_i + \mathbf{x}_j^\top \mathbf{x}_j - 2\mathbf{x}_i^\top \mathbf{x}_j$  and rearrange terms.
- Noting that  $\mathbf{X}^* = \mathbf{X} - \frac{1}{m}\mathbf{X}\mathbf{1}\mathbf{1}^\top$  and plugging into the equation  $\mathbf{K}^* = \mathbf{X}^{*\top} \mathbf{X}^*$  yields the result.
- Note that the scalar form of the equation in (b) is

$$\mathbf{K}_{ij}^* = \mathbf{K}_{ij} - \frac{1}{m} \sum_{k=1}^m \mathbf{K}_{ik} - \frac{1}{m} \sum_{k=1}^m \mathbf{K}_{kj} + \frac{1}{m^2} \sum_k \sum_l \mathbf{K}_{k,l}.$$

Substituting with the equation  $\mathbf{D}_{ij}^2 = \mathbf{K}_{ii} + \mathbf{K}_{jj} - 2\mathbf{K}_{ij}$  from (a) and simplifying yields the result.

- We first observe that  $-\frac{1}{2}\mathbf{H}\mathbf{D}\mathbf{H} = -\frac{1}{2}(\mathbf{D} - \frac{1}{m}\mathbf{D}\mathbf{1}\mathbf{1}^\top - \frac{1}{m}\mathbf{1}\mathbf{1}^\top\mathbf{D} + \frac{1}{m^2}\mathbf{1}\mathbf{1}^\top\mathbf{D}\mathbf{1}\mathbf{1}^\top)$ . By inspection, the matrix expression on the RHS corresponds to the scalar expression with four terms on the RHS of the equation in (c).

## 15.4 Nyström method

- For the first part of question, note that  $\mathbf{W}$  is SPSD if  $\mathbf{x}^\top \mathbf{W} \mathbf{x} \geq 0$  for all  $\mathbf{x} \in \mathbb{R}^l$ . This condition is equivalent to  $\mathbf{y}^\top \mathbf{K} \mathbf{y} \geq 0$  for all  $\mathbf{y} \in \mathbb{R}^m$  where  $y_i = 0$  for  $l+1 \leq i \leq m$ . Since  $\mathbf{K}$  is SPSD by assumption, this latter condition holds. For the second part, we write  $\widetilde{\mathbf{K}}$  in block form as

$$\widetilde{\mathbf{K}} = \begin{bmatrix} \mathbf{W} \\ \mathbf{K}_{21} \end{bmatrix} \mathbf{W}^\dagger \begin{bmatrix} \mathbf{W} & \mathbf{K}_{21}^\top \end{bmatrix} = \begin{bmatrix} \mathbf{W} & \mathbf{K}_{21}^\top \\ \mathbf{K}_{21} & \mathbf{K}_{21} \mathbf{W}^\dagger \mathbf{K}_{21}^\top \end{bmatrix}.$$

Comparison with the block form of  $\mathbf{K}$  then immediately yields the desired result.

- (b) Observe that  $\mathbf{C} = \mathbf{X}^\top \mathbf{X}'$  and  $\mathbf{W} = \mathbf{X}'^\top \mathbf{X}'$ . Thus,

$$\tilde{\mathbf{K}} = \mathbf{C}\mathbf{W}^\dagger \mathbf{C}^\top = \mathbf{X}^\top \mathbf{X}' (\mathbf{X}'^\top \mathbf{X}')^\dagger \mathbf{X}'^\top \mathbf{X} = \mathbf{X}^\top \mathbf{U}_{X'} \mathbf{U}_{X'}^\top \mathbf{X} = \mathbf{X}^\top \mathbf{P}_{U_{X'}} \mathbf{X}.$$

- (c) Yes. Using the expression for  $\tilde{\mathbf{K}}$  in (b) and the idempotency of orthogonal projection matrices, we can write  $\tilde{\mathbf{K}} = \mathbf{X}^\top \mathbf{P}_{U_{X'}} \mathbf{X} = \mathbf{A}^\top \mathbf{A}$ , where  $\mathbf{A} = \mathbf{P}_{U_{X'}} \mathbf{X}$ .
- (d) Since  $\mathbf{K} = \mathbf{X}^\top \mathbf{X}$ ,  $\text{rank}(\mathbf{K}) = \text{rank}(\mathbf{X}) = r$ . Similarly,  $\mathbf{W} = \mathbf{X}'^\top \mathbf{X}'$  implies  $\text{rank}(\mathbf{W}) = \text{rank}(\mathbf{X}') = r$ . The columns of  $\mathbf{X}'$  are columns of  $\mathbf{X}$ , and they thus span the columns of  $\mathbf{X}$ . Hence,  $\mathbf{U}_{X'}$  is an orthonormal basis for  $\mathbf{X}$ , i.e.,  $\mathbf{I}_N - \mathbf{P}_{U_{X'}} \in \text{Null}(\mathbf{X})$ , and by part (b) of this exercise we have  $\mathbf{K} - \tilde{\mathbf{K}} = \mathbf{X}^\top (\mathbf{I}_N - \mathbf{P}_{U_{X'}}) \mathbf{X} = \mathbf{0}$ .
- (e) Storage of  $\mathbf{K}$  requires roughly 3200 TB, i.e.,

$$(20 \times 10^6)^2 \text{ entries} \times 8 \text{ bytes/entry} \times \frac{1 \text{ TB}}{10^{12} \text{ bytes}} = 3200 \text{ TB}.$$

Storage of  $\mathbf{C}$  requires roughly 160 GB, i.e.,

$$(20 \times 10^6 \times 10^3) \text{ entries} \times 8 \text{ bytes/entry} \times \frac{1 \text{ GB}}{10^9 \text{ bytes}} = 160 \text{ GB}.$$

Note that the computed numbers do not account for the symmetry of  $\mathbf{K}$  (doing so would change the storage requirements by less than a factor of two).

## Chapter C

- C.1 For any  $\delta > 0$ , let  $t = f^{-1}(\delta)$ . Plugging this in  $\mathbb{P}[X > t] \leq f(t)$  yields  $\mathbb{P}[X > f^{-1}(\delta)] \leq \delta$ , that is  $\mathbb{P}[X \leq f^{-1}(\delta)] \geq 1 - \delta$ .

- C.2 By definition of expectation and using the hint, we can write

$$\mathbb{E}[X] = \sum_{n \geq 0} n \mathbb{P}[X = n] = \sum_{n \geq 1} n(\mathbb{P}[X \geq n] - \mathbb{P}[X \geq n+1]).$$

Note that in this sum, for  $n \geq 1$ ,  $\mathbb{P}[X \geq n]$  is added  $n$  times and subtracted  $n-1$  times, thus  $\mathbb{E}[X] = \sum_{n \geq 1} \mathbb{P}[X \geq n]$ .

More generally, by definition of the Lebesgue integral, for any non-negative random variable  $X$ , the following identity holds:

$$\mathbb{E}[X] = \int_0^{+\infty} \mathbb{P}[X \geq t] dt.$$

## Chapter D

- D.2 Estimating label bias. Let  $\hat{p}_+$  be the fraction of positively labeled points in  $\mathcal{S} = (x_1, \dots, x_m)$ :

$$\hat{p}_+ = \frac{1}{m} \sum_{i=1}^m 1_{f(x_i)=+1}$$

Since the points are drawn i.i.d.,

$$\mathbb{E}[\hat{p}_+] = \frac{1}{m} \sum_{i=1}^m \mathbb{E}_{\mathcal{S} \sim \mathcal{D}^m} [1_{f(x_i)=+1}] = \mathbb{E}_{\mathcal{S} \sim \mathcal{D}^m} [1_{f(x_1)=+1}] = \mathbb{E}_{x \sim \mathcal{D}} [1_{f(x)=+1}] = p_+.$$

Thus, by Hoeffding's inequality, for any  $\epsilon > 0$ ,

$$\mathbb{P}[|p_+ - \hat{p}_+| > \epsilon] \leq 2e^{-2m\epsilon^2}.$$

Setting  $\delta$  to match the right-hand side yields the result.

- D.3 Biased coins

- (a) By definition of the error of Oskar's prediction rule,

$$\begin{aligned}
 \text{error}(f_o) &= \mathbb{P}[f_o(S) \neq x] \\
 &= \mathbb{P}[f_o(S) = x_A \wedge x = x_B] + \mathbb{P}[f_o(S) = x_B \wedge x = x_A] \\
 &= \mathbb{P}\left[N(S) < \frac{m}{2} \mid x = x_B\right] \mathbb{P}[x = x_B] + \\
 &\quad \mathbb{P}\left[N(S) \geq \frac{m}{2} \mid x = x_A\right] \mathbb{P}[x = x_A] \\
 &= \frac{1}{2} \mathbb{P}\left[N(S) < \frac{m}{2} \mid x = x_B\right] + \frac{1}{2} \mathbb{P}\left[N(S) \geq \frac{m}{2} \mid x = x_A\right] \\
 &\geq \frac{1}{2} \mathbb{P}\left[N(S) \geq \frac{m}{2} \mid x = x_A\right].
 \end{aligned}$$

- (b) Note that  $\mathbb{P}[N(S) \geq \frac{m}{2} \mid x = x_A] = \mathbb{P}[B(m, p) \geq k]$ , with  $p = 1/2 - \epsilon/2$ ,  $k = \frac{m}{2}$ , and  $mp \leq k \leq m(1-p)$ . Thus, by Slud's inequality (section D.5)

$$\text{error}(f_o) \geq \frac{1}{2} \mathbb{P}\left[N \geq \frac{m\epsilon/2}{\sqrt{1/4(1-\epsilon^2)m}}\right] = \frac{1}{2} \mathbb{P}\left[N \geq \frac{\sqrt{m}\epsilon}{\sqrt{1-\epsilon^2}}\right].$$

Using the second inequality of the appendix, we now obtain

$$\text{error}(f_o) \geq \frac{1}{4} \left(1 - \sqrt{1 - e^{-u^2}}\right),$$

with  $u = \frac{\sqrt{m}\epsilon}{\sqrt{1-\epsilon^2}}$ , which coincides with (D.29).

- (c) If  $m$  is odd, since  $\mathbb{P}\left[N(S) \geq \frac{m}{2} \mid x = x_A\right] \geq \mathbb{P}\left[N(S) \geq \frac{m+1}{2} \mid x = x_A\right]$ , we can use the lower bound

$$\text{error}(f_o) \geq \frac{1}{2} \mathbb{P}\left[N(S) \geq \frac{m+1}{2} \mid x = x_A\right].$$

Thus, in both cases we can use the lower bound expression with  $\lceil m/2 \rceil$  instead of  $m/2$ .

- (d) If  $\text{error}(f_o)$  is at most  $\delta$ , then  $\frac{1}{4} \left[1 - \left[1 - e^{-\frac{2\lceil m/2 \rceil \epsilon^2}{1-\epsilon^2}}\right]^{\frac{1}{2}}\right] < \delta$ , which gives

$$e^{-\frac{2\lceil m/2 \rceil \epsilon^2}{1-\epsilon^2}} < 1 - (1 - 4\delta)^2 = 4\delta(2 - 4\delta) = 8\delta(1 - 2\delta),$$

and

$$m > 2 \left\lceil \frac{1 - \epsilon^2}{2\epsilon^2} \log \frac{1}{8\delta(1 - 2\delta)} \right\rceil.$$

The lower bound varies as  $\frac{1}{\epsilon^2}$ .

- (e) Let  $f$  be an arbitrary rule and denote by  $F_A$  the set of samples for which  $f(S) = x_A$  and by  $F_B$  the complement. Then, by definition of the error,

$$\begin{aligned}
 \text{error}(f) &= \sum_{S \in F_A} \mathbb{P}[S \wedge x_B] + \sum_{S \in F_B} \mathbb{P}[S \wedge x_A] \\
 &= \frac{1}{2} \sum_{S \in F_A} \mathbb{P}[S|x_B] + \frac{1}{2} \sum_{S \in F_B} \mathbb{P}[S|x_A] \\
 &= \frac{1}{2} \sum_{\substack{S \in F_A \\ N(S) < m/2}} \mathbb{P}[S|x_B] + \frac{1}{2} \sum_{\substack{S \in F_A \\ N(S) \geq m/2}} \mathbb{P}[S|x_B] + \\
 &\quad \frac{1}{2} \sum_{\substack{S \in F_B \\ N(S) < m/2}} \mathbb{P}[S|x_A] + \frac{1}{2} \sum_{\substack{S \in F_B \\ N(S) \geq m/2}} \mathbb{P}[S|x_A].
 \end{aligned}$$

Now, if  $N(S) \geq m/2$ , clearly  $\mathbb{P}[S|x_B] \geq \mathbb{P}[S|x_A]$ . Similarly, if  $N(S) < m/2$ , clearly  $\mathbb{P}[S|x_A] \geq \mathbb{P}[S|x_B]$ . In view of these inequalities,  $\text{error}(f)$  can be lower bounded as

follows

$$\begin{aligned}
 error(f) &\geq \frac{1}{2} \sum_{\substack{S \in F_A \\ N(S) < m/2}} \mathbb{P}[S|x_B] + \frac{1}{2} \sum_{\substack{S \in F_A \\ N(S) \geq m/2}} \mathbb{P}[S|x_A] + \\
 &\quad \frac{1}{2} \sum_{\substack{S \in F_B \\ N(S) < m/2}} \mathbb{P}[S|x_B] + \frac{1}{2} \sum_{\substack{S \in F_B \\ N(S) \geq m/2}} \mathbb{P}[S|x_A] \\
 &= \frac{1}{2} \sum_{S: N(S) < m/2} \mathbb{P}[S|x_B] + \frac{1}{2} \sum_{S: N(S) \geq m/2} \mathbb{P}[S|x_A] \\
 &= error(f_o).
 \end{aligned}$$

Oskar's rule is known as the *maximum likelihood* solution.